

INTRODUCTION:

MODULI PROBLEMS

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"DEFINITION"

A moduli problem is a class of geometric objects which one tries to view as another geometric object.

OUR BASIC EXAMPLE will be the

"moduli problem of smooth n -pointed curves of genus g :

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For every scheme S , we put

$\mathcal{M}_{g,n}(S)$ = category with

- objects: $(X \xrightarrow{f} S, x_1, \dots, x_n)$ where

- $f: X \rightarrow S$ is smooth, proper, with geometric fibres 1-dimensional, connected, of genus g , and

- $x_1, \dots, x_n: S \rightarrow X$ are disjoint sections

- morphisms: S -isomorphisms respecting the sections.

$\underline{\mathcal{M}}_{g,n}(S)$ = the set of isomorphism classes

of objects of $\mathcal{M}_{g,n}(S)$.

The latter is just a set, while by construction $\mathcal{M}_{g,n}(S)$ is a groupoid (=category with all maps invertible).

(3)

For every morphism $S' \rightarrow S$, we have

base change functors

$$\mathcal{M}_{g,n}(S) \rightarrow \mathcal{M}_{g,n}(S')$$

whence natural maps

$$\underline{\mathcal{M}}_{g,n}(S) \rightarrow \underline{\mathcal{M}}_{g,n}(S')$$

making each $\underline{\mathcal{M}}_{g,n}$ into a functor

$$\underline{\mathcal{M}}_{g,n}: (\text{Schemes})^\circ \rightarrow (\text{Sets})$$

In general, $\underline{M}_{g,n}$ is NOT (representable by) a scheme.

(In fact, it is a scheme iff $n > 2g + 2$).

For instance, one can find two curves

$$\begin{array}{ccc} X & & X' \\ & \searrow & \swarrow \\ & S & \end{array}$$

which are not isomorphic, but locally isomorphic over S (for the étale topology, or even the Zariski topology).

Thus, $\underline{M}_{g,0}$ (any g) is not a sheaf for these topologies, hence not representable.

(And, of course, if we take the associated sheaf we lose even more information)

(5)

MAIN IDEA:

$M_{g,n}$ is a BETTER OBJECT
to look at than $\underline{M}_{g,n}$:

- it is the natural object one tries to study (no loss of information)
- it has good | local-to-global properties
| sheaf-theoretic
| descent
- it has good approximations by schemes

Of course, in a sense it is a more complicated object ($M_{g,n}(S)$ is a groupoid, not a set).

FIBERED GROUPOIDS

let C be a category

(typically: the category of schemes,
possibly over a fixed "base scheme")

A fibered groupoid M over C (C -groupoid)
consists of the following data:

- for each $U \in \text{ob } C$, a groupoid $M(U)$
- for each map $V \xrightarrow{f} U$ in C , a functor

$$f^*: M(U) \rightarrow M(V)$$

- for each composite map $W \xrightarrow{g} V \xrightarrow{f} U$,
an isomorphism

$$g^* f^* \xrightarrow{\sim} (fg)^*$$

of functors $M(U) \rightarrow M(W)$

+ compatibility with the associativity
of composition in C .

Examples :

• $\mathcal{C} = (\text{Schemes})$:

- $U \mapsto M_{g,m}(U)$
- $U \mapsto \text{cat. of all } U\text{-schemes}$
+ U -isomorphisms
- $U \mapsto \text{QCoh}(U) := \text{cat. of quasi-coherent}$
 \mathcal{O}_U -Modules (+isomorphisms)
- $U \mapsto \text{BUN}_m(U) := \text{cat. of locally free}$
($m \in \mathbb{N}$) \mathcal{O}_U -modules of rank m
(+isomorphisms)

• For any C , any presheaf on C , i.e. any functor
 $F: C^o \rightarrow (\text{Sets})$

defines a C -groupoid (denoted by F):

$F(U) :=$ the discrete category $F(U)$:

set of objects = $F(U)$

maps = identities.

MORPHISMS OF GROUPOIDS

If \mathcal{M} , \mathcal{N} are C -groupoids, a morphism

$$\underline{\Phi}: \mathcal{M} \rightarrow \mathcal{N}$$

consists of the following data :

- For each $U \in \text{ob } C$, a functor

$$\underline{\Phi}(U): \mathcal{M}(U) \rightarrow \mathcal{N}(U)$$

- For each $V \xrightarrow{f} U$ in C , consider the diagram

$$\mathcal{M}(U) \xrightarrow{\underline{\Phi}(U)} \mathcal{N}(U)$$

of functors:

$$\begin{array}{ccc} & \downarrow f^* & \\ \mathcal{M}(U) & \xrightarrow{\underline{\Phi}(U)} & \mathcal{N}(U) \\ \downarrow f^* & \xrightarrow{\underline{\Phi}(V)} & \downarrow f^* \\ \mathcal{M}(V) & \xrightarrow{\underline{\Phi}(V)} & \mathcal{N}(V) \end{array}$$

We require (as part of the data) an isomorphism

$$f^* \circ \underline{\Phi}(U) \xrightarrow{\sim} \underline{\Phi}(V) \circ f^*$$

of functors $\mathcal{M}(U) \rightarrow \mathcal{N}(V)$.

(+ compatibility with the associativity data).

Examples:

- $\mathcal{M}_{g,n} \longrightarrow \mathcal{M}_{g,n-1}$ ($n > 0$)

"forget the n -th marked point"

- $\mathcal{M}_{1,1} \xrightarrow{j} \mathbb{A}^1$ (viewed as a presheaf)

$(X \xrightarrow{\epsilon} S) \longmapsto j(X/S) \in \Gamma(S, \mathcal{O}_S) = \mathbb{A}^1(S)$
elliptic curve

- Let us define two morphisms

$$\Phi, \Psi: \mathcal{M}_{1,1} \rightarrow \text{BUN}_1$$

by $\Phi(X \xrightarrow[\varepsilon]{p} S) := p_* \Omega^1_{X/S}$

$$\Psi(X \xrightarrow[\varepsilon]{p} S) := \varepsilon^* \Omega^1_{X/S}$$

These are different morphisms, but $\Phi(X)$ and $\Psi(X)$ are known to be canonically isomorphic.

So there should be a notion of (iso)morphism between morphisms!

Easy exercise : we obtain in this way
an equivalence of categories

$$\{1\text{-morphisms } U \rightarrow N\} \xrightarrow{\cong} N(U).$$

For instance : if S is a scheme,
a 1-morphism

$$S \longrightarrow M_{g,n}$$

is "the same thing" as an object of $M_{g,n}(S)$.

PRODUCTS

Given a diagram of C -groupoids

$$\begin{array}{ccc} \mathcal{N} & & \\ \downarrow \Phi & & \\ \mathcal{M} & \xrightarrow{\Phi} & \mathcal{P} \end{array}$$

there is a "fibre product" groupoid,
 assigning to each $V \in \text{ob } C$ the
 fibre product category

$$\mathcal{M}(V) \times_{\mathcal{P}(V)} \mathcal{N}(V)$$

whose objects are triples

$$(X, Y, \alpha)$$

with $\left\{ \begin{array}{l} X \in \text{ob } \mathcal{M}(V) \\ Y \in \text{ob } \mathcal{N}(V) \\ \alpha: \Phi(X) \cong \Psi(Y) \text{ (isomorphism} \\ \text{in } \mathcal{P}(V) \end{array} \right.$

Example:

Assume: \mathcal{M} is a C -groupoid,

U_1, U_2 objects of C

$X_i \in \text{ob } \mathcal{M}(U_i) \quad (i=1, 2)$

Viewing X_i as a 1-morphism $U_i \rightarrow \mathcal{M}$,
we get a diagram:

$$\begin{array}{ccc} & U_2 & \\ & \downarrow X_2 & \\ U_1 & \xrightarrow{X_1} & \mathcal{M} \end{array}$$

What is $U_1 \times_{X_1, \mathcal{M}, X_2} U_2$?

Answer: it is the presheaf on C given by

$$T \mapsto \left\{ (u_1, u_2, \alpha) : u_1 = \text{a morphism } T \rightarrow U_1, \right.$$

$$u_2 = \text{a morphism } T \rightarrow U_2,$$

$$\left. \alpha: u_1^* X_1 \xrightarrow{\sim} u_2^* X_2 \text{ in } \mathcal{M}(T) \right\}$$

or, in standard notations,

$$\underline{\text{Isom}}_{U_1 \times U_2} (p_1^* X_1, p_2^* X_2)$$

Example:

Consider the 1-morphism "forget the marked point"

$$\Phi: \mathcal{M}_{g,1} \rightarrow \mathcal{M}_{g,0}$$

If S is a scheme and X is an S -curve of genus g , then the fibre product $S \times_{X, \mathcal{M}_{g,0}, \Phi} \mathcal{M}_{g,1}$ is the S -curve X ! In other words, we have a Cartesian diagram

$$\begin{array}{ccc} X & \longrightarrow & \mathcal{M}_{g,1} \\ \downarrow & & \downarrow \Phi \\ S & \xrightarrow{X} & \mathcal{M}_{g,0} \end{array}$$

which shows that $\mathcal{M}_{g,1}$ can be seen as the universal curve over $\mathcal{M}_{g,0}$.

REPRESENTABLE MORPHISMS

(For safety, assume \mathcal{C} has fibre products)A 1-morphism $\Phi: \mathcal{M} \rightarrow \mathcal{N}$

is representable if for each $U \in \text{ob } \mathcal{C}$
 and $X: U \rightarrow \mathcal{N}$ (i.e. object of $\mathcal{N}(U)$)
 the fibre product $U \times_{X, \mathcal{N}, \Phi} \mathcal{M}$ is a
 presheaf, representable by an object of \mathcal{C} .

For instance, the "forgetful" morphism

$$\mathcal{M}_{g,1} \rightarrow \mathcal{M}_{g,0}$$

is representable.

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Another example: for $g \geq 2$, consider

$$3K: \mathcal{M}_{g,0} \longrightarrow \text{BUN}_r \quad (r=5g-5)$$

$$(X \xrightarrow{f} U) \mapsto f_* \omega_{X/U}^{\otimes 3}$$

I claim that $3K$ is representable:

Pick a scheme T and a 1-morphism $T \rightarrow \text{BUN}_r$
(that is, a locally free sheaf \mathcal{E} on T , of rank r):

$$\begin{array}{ccc} \mathcal{N} & \longrightarrow & \mathcal{M}_{g,0} \\ \downarrow & & \downarrow 3K \\ T & \xrightarrow{\mathcal{E}} & \text{BUN}_r \end{array}$$

Then, for a T -scheme T , we have

$\mathcal{N}(T) = \text{cat. of curves } X \xrightarrow{f} T \text{ of genus } g,$
plus isomorphism $f_* \omega_{X/T}^{\otimes 3} \xrightarrow{\sim} \mathcal{E}_T$.

Such a curve is naturally (3-canonically)
embedded in $\mathbb{P}(\mathcal{E}_T)$. Putting $\mathbb{P} = \mathbb{P}(\mathcal{E})$, we
obtain an equivalence:

$\mathcal{N}(T) \xrightarrow{\sim} \text{cat. of embedded smooth curves of}$
genus $g : X \hookrightarrow \mathbb{P} \times_T T,$
plus isomorphism $\mathcal{O}(1)|_X \simeq \omega_{X/T}^{\otimes 3}$

The representability then follows from Hilbert scheme
theory.

PROPERTIES OF REPRESENTABLE MORPHISMS

If P is a property (i.e. a class) of morphisms of \mathcal{C} , which is stable by base change, it makes sense to say that a representable 1-morphism $\Phi: M \rightarrow N$ has property P .

For instance, in the above examples,

$\Phi: M_{g,1} \rightarrow M_{g,0}$ is proper and smooth.

$3K: M_{g,0} \rightarrow \text{BUN}_{S_{g-5}}$ is surjective and smooth.

USING THE DIAGONAL

(We assume C has fibre products)

Proposition For a C -groupoid \mathcal{M} , the following conditions are equivalent:

(i) The diagonal 1-morphism

$$\Delta_{\mathcal{M}} : \mathcal{M} \longrightarrow \mathcal{M} \times \mathcal{M}$$

$$x \longmapsto (x, x)$$

is representable.

(ii) For all $U \in \text{ob } C$ and $x, y \in \text{ob } \mathcal{M}(U)$, the presheaf $\underline{\text{Isom}}_{\mathcal{M}}(x, y)$ is representable (by an object of C/U).

(iii) For each $U \in \text{ob } C$, every 1-morphism $U \rightarrow \mathcal{M}$ is representable.

WARNING: $\Delta_{\mathcal{M}}$ is NOT in general a "monomorphism", i.e. a fully faithful functor.

In fact:

$\Delta_{\mathcal{M}}(U) : \mathcal{M}(U) \rightarrow (\mathcal{M} \times \mathcal{M})(U)$ is fully faithful for each U



\mathcal{M} is (associated to) a presheaf on C

Note : The properties in the above Proposition are satisfied for $M = M_{g,n}$.

In fact,

$$\Delta_{M_{g,n}} : M_{g,n} \longrightarrow M_{g,n} \times M_{g,n}$$

is representable, separated, of finite type.

If $2g - 2 + n > 0$, it is finite unramified
(objects of $M_{g,n}$ have no infinitesimal automorphisms)

If $n > 2g + 2$ then it is a monomorphism,
in fact a closed immersion

(objects of $M_{g,n}$ have no nontrivial automorphisms,
and $M_{g,n}$ is a presheaf in this case)

But for instance, for $g = n = 0$, consider
 $X : \text{Spec } \mathbb{Z} \rightarrow M_{0,0}$ defined by \mathbb{P}^1 : then

Then $\text{Spec } \mathbb{Z} \times_{X, M_{0,0}, X} \text{Spec } \mathbb{Z} \simeq \underline{\text{PGL}}_{2, \mathbb{Z}}$