# Set Cascades and Vector Valleys in Pascal's Triangle

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#### Abstract

Let N be the positive integers. Let T(n) be the set of all  $F \subseteq N$  of size |F| = n. The shadow  $\Delta F$  of F is the subset of T(n-1) obtained by deleting an element of F in all n possible ways. If  $R \subseteq T(n)$  then  $\Delta R = \bigcup \{F \in R\} \Delta F$ . If we write |R| as an n-Cascade we can immediately write down the best possible lower bound for  $|\Delta R|$ .

Let V(n) be the set of all vectors of dimension n with coordinates 0 or 1. The shadow  $\Delta A$  of  $A \in V(n)$  is the subset of V(n-1) obtained by deleting a coordinate of A in all ways. If  $W \subseteq V(n)$  we write |W| as an n-Valley and get the best possible lower bound for  $|\Delta W|$ .

### 1 Shadows of sets.

Let N be the set of all positive integers. For  $n \in N$  let T(n) be the set of all  $F \subseteq N$  of size |F| = n.

Ex.1. We have  $T(1) = \{1, 2, 3, ...\}$  and  $T(2) = \{12, 13, 23, 14, 24, 34, 15, ...\}$  and  $T(3) = \{123, 124, 134, 234, 125, ...\}$ .

If  $F \in T(n)$  the shadow  $\Delta F$  of F is the  $\Delta F \subseteq T(n-1)$  obtained by deleting an element from F in all n possible ways, so  $|\Delta F| = n$ .

Ex.2. If F = 123 then  $\Delta F$  is  $\{12, 13, 23\}$ .

If  $R \subseteq T(n)$  then  $\Delta R = \bigcup \{F \in R\} \Delta F$ .

Ex.3. Let  $P = \{12, 13, 23, 14, 24, 34\}$  and  $Q = \{12, 13, 23, 24, 25\}$ . Let  $R_1 = \{12, 13, 23, 14, 24, 34\}$ 

 $\{123, 124, 134\}, R_2 = \{123, 124, 234\}, R_3 = \{123, 124, 125\} \text{ and } R_4 = \{123, 246, 256\}.$ Then  $\Delta R_1 = \Delta R_2 = P$  and  $\Delta R_3 = Q \cup \{14, 15\}$  and  $\Delta R_4 = Q \cup \{26, 46, 56\}.$ 

If  $R \subseteq T(n)$  then clearly  $|\Delta R| \le n|R|$  but we want to know how small  $|\Delta R|$  can be in terms of |R|. To answer we order each T(n). We do this by the rule F < G, in words F is before G, if  $max\{F \setminus G\} < max\{G \setminus F\}$ .

Ex.4. Since  $max\{124 \setminus 345\} = max\{12\} = 2 < 5 = max\{35\} = max\{345 \setminus 124\}$  we have 124 < 345 in T(3).

We see this order in Ex.1,3. By an *initial section IS* we mean the first so many sets in order of T(n). Thus P,  $R_1$  are IS in Ex.3. An important fact is that the shadow of an IS is an IS.

**Theorem 1** (Kruskal 1963, Katona 1966). Let  $R \subseteq T(n)$ . If S is the IS of T(n) with |R| = |S| then  $|\Delta R| \ge |\Delta S|$ .

This theorem has been proved in various ways. Here we will explain the idea of the proof in [3].

#### 2 Shifts of sets.

Given  $R \subseteq T(n)$  we want to shift it step by step towards the beginning of the order, until it becomes an IS. Each shift must not increase  $|\Delta R|$  or change |R|, and Theorem 1 will be proved. If  $i, j \in N$  and i < j then the shift  $i \leftarrow j$  changes a j into an i in R wherever possible. So for  $F \in R$  the shift  $(i \leftarrow j)F$  is  $(F \setminus j) \cup i$  if  $i \notin F$  and  $j \in F$  and  $(F \setminus j) \cup i \notin R$ , otherwise it is F. Of course  $(i \leftarrow j)R = \{(i \leftarrow j)F : F \in R\}$ . In Ex.3 we have  $R_1 = (1 \leftarrow 2)R_2$  and  $R_3 = (1 \leftarrow 6)R_4$  with  $7 = |\Delta R_3| < |\Delta R_4| = 8$ . In fact  $\Delta(1 \leftarrow 6)R_4 = \Delta R_3 \subseteq (1 \leftarrow 6)\Delta R_4$  which is a special case of

### Lemma 1 $Shadow(shiftR) \subseteq Shift(shadowR)$ .

To prove this for  $i \leftarrow j$  you consider an arbitrary  $G \in \Delta(i \leftarrow j)R$ . There is an  $h \in N$  with  $G \cup h \in (i \leftarrow j)R$ . So there is an  $F \in R$  with  $G \cup h = (i \leftarrow j)F$ . Suppose  $j \notin F$ . Then  $G \cup h = F$  so  $j \notin G \in \Delta R$  and  $G \in (i \leftarrow j)\Delta R$  as required. An interested reader can easily sort out the other cases.

Now in Ex.3 the IS is  $R_1$  but we cannot get it from  $R_3$  by  $i \leftarrow j$  shifts. We need more general shifts. We take sets  $I, J \subseteq N$  with |I| = |J| and I < J and  $I \cap J$  empty. The shift  $I \leftarrow J$  of  $F \in R$  is  $(F \setminus J) \cup I$  if  $I \cap F$  is empty,  $J \subseteq F$  and  $(F \setminus J) \cup I \not\in R$ , otherwise it is F. We have  $R_2 = (34 \leftarrow 15)R_3$  in Ex.3. If R is not an IS there exist such I, J with |I| minimal such that  $(I \leftarrow J)R \neq R$ . For this shift Lemma 1 still holds, but it is a little harder to write out the proof. We trust that the reader trusts Lemma 1 and Theorem 1.

## 3 Cascade Representations of Integers.

To get numbers out of Theorem 1 we let S be an arbitrary IS of T(n). There is a largest  $m_n \in N$  such that S contains

all 
$$\binom{m_n}{n}$$
 sets  $F$  with  $F \subseteq \{1, ..., m_n\}, |F| = n$ .

If there are more sets in S they must possess the number  $m_n + 1$ . Also there must be a largest  $m_{n-1}$  so that S contains

all 
$$\binom{m_{n-1}}{n-1}$$
 sets  $G \cup \{m_n+1\}$  with  $G \subseteq \{1,...,m_{n-1}\}, |G| = n-1$ .

If there are more sets in S there is a largest  $m_{n-2}$  so that S contains

all 
$$\binom{m_{n-2}}{n-2}$$
 sets  $H \cup \{m_{n-1}+1, m_n+1\}$  with  $H \subseteq \{1, ..., m_{n-2}\}, |H| = n-2,$ 

and so on. Assume we end with  $m_e$  and put

$$k = {\binom{m_n}{n}} + {\binom{m_{n-1}}{n-1}} + \dots + {\binom{m_e}{e}} \text{ where } m_n > m_{n-1} > \dots > m_e \ge e \ge 1,$$
 (1)

so k = |S|. We call (1) an n-Cascade for k. If we plot (1) on Figure 1, the successive binomial coefficients are in successive columns, each not below the one before. So the plot resembles a cascade of water getting deeper as it runs.

Figure 1. Two orientations of Pascal's Triangle.

From the above construction we see that  $\Delta S$  contains

all 
$$\binom{m_n}{n-1}$$
 sets  $G$  with  $G \subseteq \{1, ..., m_n\}, |G| = n-1,$ 

all 
$$\binom{m_{n-1}}{n-2}$$
 sets  $H \cup \{m_n+1\}$  with  $H \subseteq \{1,..,m_{n-1}\}$ ,  $|H| = n-2$ ,

and so on. Thus we just move the n-Cascade (1), using the rule in Figure 1, to get the (n-1)-Cascade for  $|\Delta S|$ . There may be one other  $m_1, ..., m_e$  which satisfies (1), it has e = 0. We can use either. Also because  $\Delta S$  is an IS we can repeat to get  $|\Delta \Delta S|$ ...

Ex.5. (Lehmer 1964). The 9167-th in S(12) is 1, 2, 4, 5, 8, 10, 11, 12, 13, 15, 16, 18.

Ex.6. (Clements 1974). If  $P, Q, S \subseteq T(n)$  and S is the IS with |P| + |Q| = |S| then  $|\Delta P| + |\Delta Q| \ge |\Delta S|$ . Just add a large h to every integer of Q.

Ex.7. (Hilton 1979). let P, Q, R be IS with  $P, Q \subseteq T(n)$ ,  $R \subseteq T(n-1)$  and |P| = |Q| + |R| then  $|\Delta P| \le |\Delta R| + max\{|R|, |\Delta Q|\}$ .

#### 4 The V-order for vectors.

For  $n \in N$  let V(n) be all n-dimensional 0, 1 vectors.

Ex.8.  $V(2) = \{00, 10, 01, 11\}$  and  $V(3) = \{000, 100, 010, 001, 110, 101, 011, 111\}$ . If  $A \in V(n)$  the shadow  $\Delta A \subseteq V(n-1)$  is obtained by deleting a coordinate of A in all ways.

Ex.9. So  $\Delta(011100) = \{01100, 11100, 01110\}.$ 

Also  $\Delta W = \bigcup \{A \in W\} \Delta A$  for  $W \subseteq V(n)$ . We need some notation. For  $A = (a_1, ..., a_n) \neq B = (b_1, ..., b_n)$  in V(n) we put  $\nu A = a_1 + ... + a_n$  and  $\alpha(A, B)$  is the first i with  $a_i \neq b_i$ . Also " $A = (a_1, ..., a_{n-1})$  and A" =  $a_n$  and  $dA = (d, a_1, ..., a_n)$  for d = 0, 1.

**Definition 1** (V-order). If  $A, B \in V(n)$  then A < B if either (i)  $\nu A < \nu B$ , or (ii)  $\nu A = \nu B$  and  $1 = a_j > b_j = 0$  with  $j = \alpha(A, B)$ .

This order is seen in Ex.8,9. Using this we proved in [2] the main result, which is Theorem 2 below. We will not attempt the proof of Theorem 2 here. It seems of necessity to be hard due to the fact there are infinitely many other orders equally good for it.

**Theorem 2** Let  $W \subseteq V(n)$ . If S is the IS of V(n) with |W| = |S| then  $|\Delta W| \ge |\Delta S|$ .

### 5 The sequence of free sequences.

Using V-order let  $V(n,t)=\{A\in V(n): \nu A=t\}=A_1< A_2< ...< A_e$  so  $e=\binom{n}{t}$ . Also let  $\Psi(n,t)$  be the 0,1 sequence  $A_1$ ",  $A_2$ ", ...,  $A_e$ ". Here  $V,\Psi$  are empty unless  $0\leq t\leq n$ . We see that V(2,0) is 00, and V(2,1) is 10<01, and V(2,2) is 11. In general V(n+1,t) is 1V(n,t-1) followed by 0V(n,t), so  $\Psi(n+1,t)$  is  $\Psi(n,t-1)$  followed by  $\Psi(n,t)$ , and we can contruct  $\Psi$  without using V-order. The n-th free sequence  $\Psi_n$  has length  $2^n$  and is  $\Psi(n,0),\Psi(n,1),...,\Psi(n,n)$ .

Let  $A \in V(n)$  and S be the IS ending A. Next put  $R = S \setminus A$  so R is the IS with |R| = |S| - 1. If  $A^n = 1$  then  $\Delta R = \Delta S$  and we say that A is free (over R), otherwise  $A^n = 0$  and  $1 + |\Delta R| = |\Delta S|$  because ("A)  $\cup$  ( $\Delta R$ ) =  $\Delta S$ . From Theorem 2 it is clear that  $|\Delta S| = |S| - c$ , where c is the sum of the first |S| terms of  $\Psi_n$ .

### 6 Valley Representations of Integers.

On Figure 1 please plot the binominal coefficients of

$$f = \begin{pmatrix} 9 \\ 0 \end{pmatrix} + \begin{pmatrix} 9 \\ 1 \end{pmatrix} + \begin{pmatrix} 9 \\ 2 \end{pmatrix} + \begin{pmatrix} 9 \\ 3 \end{pmatrix} \quad and \quad g = \begin{pmatrix} 8 \\ 3 \end{pmatrix} + \begin{pmatrix} 6 \\ 2 \end{pmatrix} + \begin{pmatrix} 5 \\ 2 \end{pmatrix} + \begin{pmatrix} 2 \\ 0 \end{pmatrix} + \begin{pmatrix} 1 \\ 0 \end{pmatrix}.$$

We think of f+g as a valley, with the 45 degree slope of f as the left side, and g as the right side. Notice that g is like a Cascade because successive coefficients are to the right, but not below, the last. It is easy to see that each k in  $0 \le k \le 2^n$  has one or two n-Valley representations.

Next we find the (f+g)-th vector A in V(9). For i=0,1,2,3 the i-th term in f corresponds to V(9,i), so  $\nu A=4$ . The number of vectors in V(9,4) which start 1 or 011 or 010111 or B=01001101 are respectively  $\binom{8}{3}$  or  $\binom{6}{2}$  or  $\binom{5}{2}$  or  $\binom{2}{0}$  or  $\binom{1}{0}$ , hence A=B0.

If S is an IS of V(n) we find the n-Valley of |S|, then we move the valley by the rule in Figure 1 to get the (n-1)-Valley of  $|\Delta S|$ . Again  $\Delta S$  is an IS and we repeat to get  $|\Delta \Delta S|$ ,...

#### 7 The T-order for Vectors.

First let p and q be two sequences  $p_1, ..., p_e$  and  $q_1, ..., q_e$  with  $p_i, q_i \in N$ . We order p < q if the least j with  $p_j \neq q_j$  has  $p_j < q_j$ .

Second we find the *type* of a vector by replacing 00 by 0 and 11 by 1 as many times as possible.

Ex.10. If A = 001111011000 then type(A) = 01010.

We order types as follows  $0 < 1 < 10 < 01 < 010 < 101 < 1010 < 0101 < \dots$  Two vectors of different types are ordered according to their types. It remains to order vectors of the same type.

Case A'' = 0. Let  $z_1$  be the number of zeros at the end of A, then  $n_1$  be the number of ones next to those zeros, and so on. The sequence  $\sigma A$  of A is  $n_1, n_2, ..., z_2, z_1$ . In Ex.10 we see A looking like  $z_3, n_2, z_2, n_1, z_1 = 2, 4, 1, 2, 3$  but  $\sigma A$  is 2, 4, 2, 1, 3. We order vectors of the same type according to the order of their sequences.

Case A'' = 1. We exchange 0 and 1 to get the complements of vectors, then we order them according to the order of their complements.

Having defined T-order we can now point out that the idea behind it is that the vectors in  $\Delta A$  have the same type as A or earlier. It turns out that the shadow of an IS is an IS, and  $A \in V(n)$  is free iff  $a_{n-1} \neq a_n$ . Most importantly Theorem 2 holds for V-order and for T-order.

### 8 Open Problems.

1) What is smallest  $R \subseteq V(n)$  with  $\Delta R = V(n-1)$ ?

2) How big can shadows of families of vectors be? More precisely for  $1 \le k \le 2^n$ 

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