

# Substitution Tilings and Substitution Tiling Spaces

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# Outline

## 1 The chair tiling

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- 4 Inverse limits

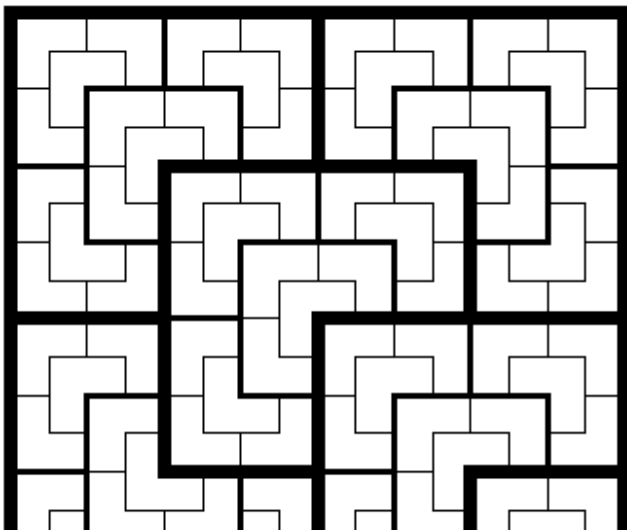
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- 5 Example calculations

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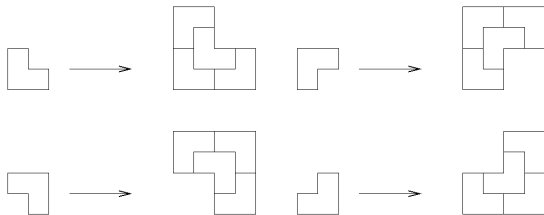
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## A big patch

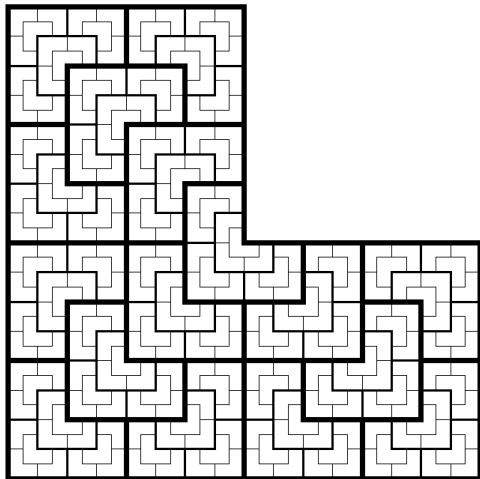




# Musical chairs



## A smaller patch



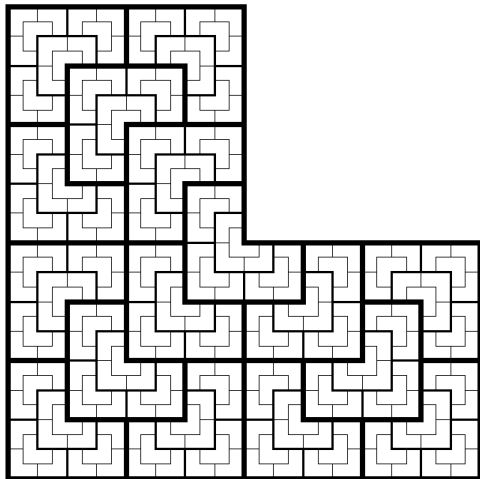
# What are $n$ -supertiles?

- Traditional view:  $\sigma^n(t) = \sigma(\sigma^{n-1}(t))$  is what you get when you apply the substitution to every tile in an  $n - 1$ -supertile.

# What are $n$ -supertiles?

- Traditional view:  $\sigma^n(t) = \sigma(\sigma^{n-1}(t))$  is what you get when you apply the substitution to every tile in an  $n - 1$ -supertile.
- Fusion perspective:  $\sigma^n(t) = \sigma^{n-1}(\sigma(t))$  is what you get when you assemble several  $n - 1$ -supertiles according to the pattern of  $\sigma(t)$ .

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## Counting chairs

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- $4 \times 4$  ways to make a 2-supertile.
- $4^n$  ways to make an  $n$ -supertile.
- Uncountably many ways to make a tiling.
- Only countably many of those are translates of a given tiling, so there are uncountably many chair tilings, up to translation.

# How different are different chair tilings?

Let  $T_1$  and  $T_2$  be different chair tilings.

- Every patch  $P$  of  $T_1$  lives in an  $n$ -supertile  $S_n$ .

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- $T_1$  and  $T_2$  have exactly the same local patterns. Same “LI class”.

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## Live is easier without shape

- In 1D, the geometry takes care of itself.
- Only need to specify the combinatorics.
- E.g., Fibonacci:  $\sigma(a) = ab$ ,  $\sigma(b) = a$ .
- Up to overall scale, tile lengths are determined by substitution matrix.

## The substitution matrix

- $M_{ij}$  counts the number of  $t_i$  tiles in  $\sigma(t_j)$ .
- For Fibonacci,  $M = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$ .



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- For period-doubling,  $\sigma(a) = ab$ ,  $\sigma(b) = aa$ ,  $M = \begin{pmatrix} 1 & 2 \\ 1 & 0 \end{pmatrix}$ .
- Don't confuse  $M$  with  $M^T$ !

## Populations and eigenvectors

- $j$ -th column of  $M$  gives population of 1-supertile  $\sigma(t_j)$ .
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- Population scales as  $\lambda^n$ , where  $\lambda$  is Perron-Frobenius eigenvalue of  $M$
- Relative density of different tiles is given by right-eigenvector.

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- $LM = \lambda L$ .
- Relative length of tiles is given by leading left-eigenvector of  $M$ .

## Example: Fibonacci

- $\sigma(a) = ab, \sigma(b) = a$
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- Leading left-eigenvector  $(\phi, 1)$ . “a” tiles are  $\phi$  times longer than “b” tiles.

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- Leading left-eigenvector  $(1, 1)$ . Both tiles have the same length.

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- Period-doubling is non-periodic because it has structures at all scales  $2^n$
- ... *abaaabababaaabaa* ...
- ... *XbXaXbXbXbXaXbXa* ...
- ... *XYXaXYXbXYXaYXYa* ...
- ... *XYXZXYXbXYXZXYXa* ...

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## Tiling metric

Define metric on set of all tilings with given tile set.

- $d(T, T') \leq \epsilon$  if  $T$  and  $T'$  agree on  $B_{1/\epsilon}(0)$ , up to  $\epsilon$ -translation.



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- Metric depends on choice of origin, but topology doesn't.
- $\lim d(T_i, T_\infty) = 0$  means  $T_i \cap K \Rightarrow T_\infty \cap K$  for all compact  $K \subset \mathbb{R}^n$ .

## Group action

$G = \mathbb{R}^n$  acts on tilings.  $T - x$  is what you get by translating all the tiles in  $T$  by  $-x$ . (Equivalently, moving the origin in  $T$  by  $+x$ .)

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A *tiling space* is a set  $\Omega$  of tilings on a fixed tile set with fixed adjacency rules such that

- $\Omega$  is translation invariant. (If  $T \in \Omega$  and  $x \in \mathbb{R}^n$ , then  $T - x \in \Omega$ .)
- $\Omega$  is closed in the tiling metric.

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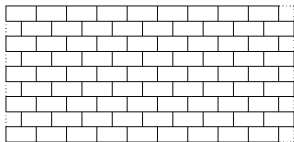


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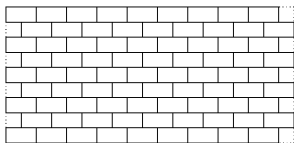
- Start with an FLC tiling  $T$ .
- Consider the set  $\{T - x\}$  of translates of  $T$ .
- $\Omega_T = \overline{\{T - x\}}$ .
- Orbit closure of  $T =$  Tiling space of  $T =$  Continuous hull of  $T$ .

# Hulls of periodic tilings

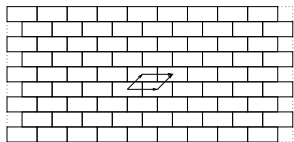


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# Hulls of periodic tilings



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A torus!

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- Hull = slinky! Connected but not path-connected.



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- Fill out near  $\infty$ . Discrete choices.
- Neighborhood  $\sim B_\epsilon \times C$ .

## Description of $\Omega_T$

### Theorem

*A tiling  $T'$  is in  $\Omega_T$  if and only if every patch of  $T'$  is found somewhere in  $T$ .*

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## Dynamical vs. combinatorial properties

- If  $T$  has FLC,  $\Omega_T$  is compact.
- If  $T$  has uniform patch frequencies,  $\Omega_T$  is uniquely ergodic.
- If  $T$  is repetitive,  $\Omega_T$  is minimal.
- Path components of  $\Omega_T$  are orbits. Typically uncountably many.

# Recognizability

## Theorem (Mossé, Solomyak)

*Let  $\sigma$  be a primitive substitution, and let  $T$  be a tiling build from  $\sigma$ . The substitution induces a surjective map  $\sigma : \Omega_T \rightarrow \Omega_T$ . This map is injective if and only if  $T$  is non-periodic.*

If  $\sigma$  is invertible, we say that the substitution is *recognizable*, or that it has the *unique decomposition property*.

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## Inverse limits in general

If  $X_0, X_1, \dots$  are spaces and  $\rho_n : X_n \rightarrow X_{n-1}$  are continuous maps,

$$X = \varprojlim X_i := \{(x_0, x_1, \dots) \in \prod X_n \mid \rho_n(x_n) = x_{n-1} \forall n\}.$$

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$X$  has the product topology.  $(x_0, x_1, \dots)$  is close to  $(y_0, y_1, \dots)$  if  $x_i \approx y_i$  for all  $i \leq N$ . I.e. if  $x_N \approx y_N$ .

# Dyadic Solenoid

Example of inverse limit space. Take

- $X_n = \mathbb{R}/(2^n\mathbb{Z}) \simeq S^1$ .

## Dyadic Solenoid

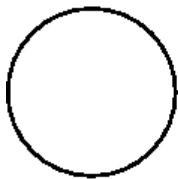
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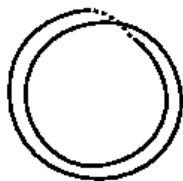
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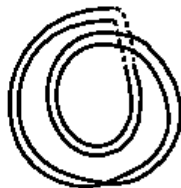
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$\Gamma_0$



$\Gamma_1$



$\Gamma_2$

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- $\varprojlim \Gamma_n =$  consistent instructions for tiling bigger and bigger regions, i.e. instructions for a complete tiling.
- So how do instructions for partial tilings turn into a CW complex?!

# Anderson-Putnam Complex

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Identify!
- $\Gamma_n = \coprod \sigma^n(t_i) / \sim$  is the **Anderson-Putnam** complex.
- $\Gamma_n$  looks just like  $\Gamma_0$ , only a factor of  $\lambda^n$  bigger.

# Anderson-Putnam Theorem

## Theorem

If  $\sigma$  is a primitive and recognizable substitution that “forces the border”, then

- $\Omega_\sigma$  is homeomorphic to  $\varprojlim(\Gamma_0, \sigma)$ .
- $\check{H}^k(\Omega_\sigma) = \varinjlim(H^k(\Gamma_0), \sigma^*)$ .

## Living without border forcing

### Theorem (Anderson-Putnam)

*If  $\sigma$  doesn't force the border, then we can get an equivalent substitution that does force the border by "collaring".*

### Theorem (Barge-Diamond)

*If a 1D substitution meets some conditions that are a lot weaker than forcing the border, then  $\Omega_\sigma$  has the same cohomology as  $\varprojlim(\Gamma_0, \sigma)$ , even though they aren't homeomorphic.*

The BD theorem applies to Fibonacci and period-doubling.



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# Dyadic solenoid

$H^1(S^1) = \mathbb{Z}$ . Substitution map induces multiplication by 2.  
 $\check{H}^1(S) = \mathbb{Z}[1/2]$

# Fibonacci

[Switch to doc-cam]

## Fibonacci results

- $\Gamma_0$  is the wedge of 2 circles, obtained by identifying all of the endpoints of the  $a$  and  $b$  tiles.
- $H^1(\Gamma_0) = \mathbb{Z}^2$
- $\sigma^*$  expressed by invertible matrix  $M^T = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$ .
- $\check{H}^1(\Omega_{Fib}) = \mathbb{Z}^2$ .