

Complex Braid Groups

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Part 2 : Standard monoids
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Notations for complex reflection groups

Let $W < GL(V)$ be a complex reflection group, $n = \dim V$

$$W = \langle \mathcal{R} \rangle \quad \mathcal{R} = \{s \in W; \dim \text{Ker}(s - 1) = n - 1\}$$

The collection of its *reflecting hyperplanes* is the *hyperplane arrangement*

$$\mathcal{A} = \{\text{Ker}(s - 1), s \in \mathcal{R}\}$$

For $H \in \mathcal{A}$, $W_H = \{w \in W; w|_H = \text{Id}_H\}$ is cyclic, isomorphic to its image under $\det : W_H \rightarrow \mathbb{C}^\times$.

The generator of W_H mapped to $\exp(2\pi i/|W_H|)$ is a reflection s_H called the *distinguished reflection* associated to H . The collection of all distinguished reflections is denoted \mathcal{R}^* .

\mathcal{R}^* is in 1-1 correspondence with \mathcal{A} ,

$$s \mapsto \text{Ker}(s - 1), \quad H \mapsto s_H$$



Classification of irreducible CRG's

The main series is made of the groups $W = G(de, e, n)$ of

- $n \times n$ monomial matrices
- with nonzero entries inside μ_r , $r = de$
- whose product belongs to μ_d .

Of course $G(r, r, n) < G(de, e, n) < G(r, 1, n)$.

- W contains **diagonal reflections**, of the form $\text{diag}(1, \dots, 1, \zeta, 1, \dots)$ if and only if $d > 1$.
- its **non-diagonal** reflections belong to $G(r, r, n) < W$ and have the form

$$\text{Id}_u \oplus \begin{pmatrix} 0 & \zeta_e^{-k} \\ \zeta_e^k & 0 \end{pmatrix} \oplus \text{Id}_{n-2-u}$$

In addition to these, there are 34 exceptional groups G_4, \dots, G_{37} , half of them in rank 2.

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Their braid groups are

- the braid group B_n for $G(1, 1, n)$
- the punctured braid group B_n^* for $G(r, 1, n) = G(d, 1, n)$ when $d > 1$
- a finite index normal subgroup $B_n^*(e)$ of B_n^* when $d > 1$ and $e > 1$
- a quotient $B_n(e)$ of $B_n^*(e)$ for $G(e, e, n) = G(r, r, n)$.

Preliminaries : Garside monoids

- A monoid is called *cancellative* if, for all $a, b, c \in M$, $ac = bc$ implies $a = b$ and $ca = cb$ implies $a = b$
- An element $a \in M$ left-divides $c \in M$ if $\exists b \in M$ $ab = c$. Then, c is a right-multiple of a , and one writes $a \prec c$. Similarly, a right-divides $c \in M$ if $\exists b \in M$ $ba = c$ and c is then a left-multiple of a , and one writes $c \succ a$.
- Two elements a, b admit a right lowest common multiple (lcm) if they admit a right common multiple $c = \text{lcm}_R(a, b)$ such that, $\forall m \in M$ $a \prec m, b \prec m \Rightarrow c \prec m$. They admit a left lcm if they admit a left common multiple $c = \text{lcm}_L(a, b)$ such that $\forall m \in M$ $m \succ a, m \succ b \Rightarrow m \succ c$.
- Two elements a, b admit a left greatest common divisor (gcd) if they admit a left common divisor $c = \text{gcd}_L(a, b)$ such that, $\forall m \in M$ $m \prec a, m \prec b \Rightarrow m \prec c$. They admit a right gcd if they admit a right common divisor $c = \text{gcd}_R(a, b)$ such that $\forall m \in M$ $a \succ m, b \succ m \Rightarrow c \succ m$.

If $M^\times = 1$ and M is cancellable, these lcm's and gcd's are uniquely defined.

An element $a \in M$ is called *reducible* if there exists $b, c \in M$ with $b, c \notin M^\times$ such that $a = bc$. It is called *irreducible* if it is not invertible and not reducible.

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An *homogeneous monoid* is a monoid M together with a *length function*, that is a monoid morphism $\ell : M \rightarrow \mathbb{N}$, such that M is generated by the elements of length > 0 .

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Under these conditions, every element of M is a product of irreducible elements. They are called the *atoms* of the monoid M .

Definition

The *group of fractions* of M is by definition a group $\text{Frac}(M)$ together with a morphism of monoids $M \rightarrow \text{Frac}(M)$ such that every $M \rightarrow G$ for G a group factors

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For S a set of generators and R a collection of relations, if M is presented as $\langle S \mid R \rangle^+$, then $\langle S \mid R \rangle$ is a presentation of $\text{Frac}(M)$.

Definition

An element of a monoid M is said to be *balanced* if the sets of its left and right divisors are the same.

Definition

An homogeneous monoid M is said to have the Garside property, or to be a Garside monoid, if it is cancellable, and if it has the following properties

- any two elements of M admit *gcd*'s and *lcm*'s on the right and on the left
- M admits a balanced element Δ whose set of divisors is finite and generates M .

The chosen element Δ is called a *Garside element* for M .

Preferred Garside element

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For $m \in M$, we set

$$\text{Div}_L(m) = \{a \in M \mid a \prec m\} \quad \text{Div}_R(m) = \{a \in M \mid m \succ a\}$$

and, if m is balanced $\text{Div}(m) = \text{Div}_L(m) = \text{Div}_R(m)$.

Complex braid groups and Garside groups

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How to deal with $\mathcal{B}_n^*(e)$?

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Finite index normal subgroups of Garside groups

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We have $C_H(x) = C_G(x) \cap H = \text{Ker}(\Phi_{C_G(x)} : C_G(x) \twoheadrightarrow F)$.

Since F is finite, from a finite set of generators of $C_G(x)$ one gets a finite set of generators of $\text{Ker}(\Phi_{C_G(x)})$ by Schreier's Lemma.

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So : is there $b \in H$ with $bc^{-1} \in C_G(x) = \langle g_1, \dots, g_r \rangle$?

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So : is there $b \in H$ with $bc^{-1} \in C_G(x) = \langle g_1, \dots, g_r \rangle$? Actually equivalent to checking whether $\Phi(c) \in \langle \Phi(g_1), \dots, \Phi(g_r) \rangle < F$.

Preliminaries : Garside interval monoids

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- $a \prec b$ means $l_S(b) = l_S(a) + l_S(a^{-1} \cdot b)$,
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For $c \in W$,

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By definition, the interval monoid attached to (W, S, c) with c balanced is defined by generators and relations

$$M(c) = \langle \text{Div}(c) \mid z = xy \text{ if } z = x \cdot y \text{ and } l_S(z) = l_S(x) + l_S(y) \rangle^+$$

Theorem

If $(\text{Div}(c), \prec)$ and $(\text{Div}(c), \succ)$ are lattices, then $M(c)$ is Garside with Garside element c .

Moreover, the poset structures on $\text{Div}(c)$ are the same inside W and inside $M(c)$.

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and both of them are Garside, with a preferred Garside element.

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Standard monoid for $G(e, e, n)$

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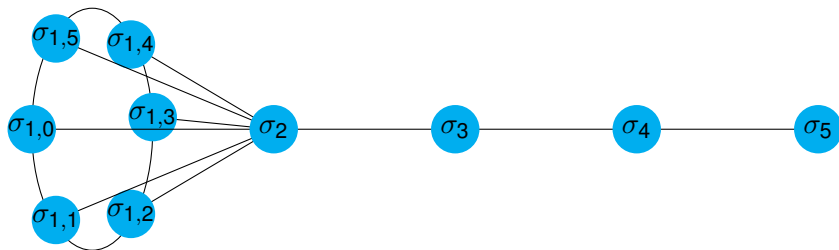
The presentation for B^+ :

$$\left\langle \begin{array}{l} \sigma_{1,k}, k \in \mathbb{Z}/e\mathbb{Z} \\ \sigma_2, \dots, \sigma_{n-1} \end{array} \middle| \begin{array}{l} \sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1} \\ \sigma_{1,k} \sigma_2 \sigma_{1,k} = \sigma_2 \sigma_{1,k} \sigma_2 \\ \sigma_{1,k} \sigma_j = \sigma_j \sigma_{1,k}, j \geq 3 \\ \sigma_{1,k} \sigma_{1,k+1} = \sigma_{1,k+1} \sigma_{1,k+2}, k \in \mathbb{Z}/e\mathbb{Z} \end{array} \right\rangle^+$$

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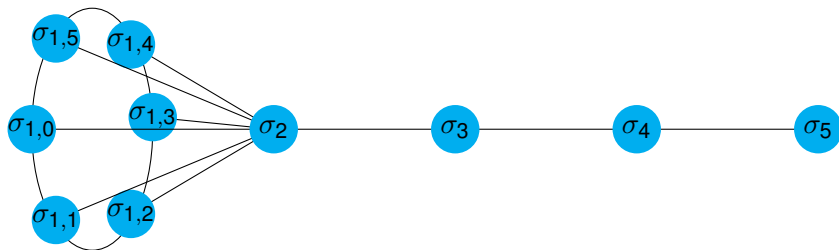
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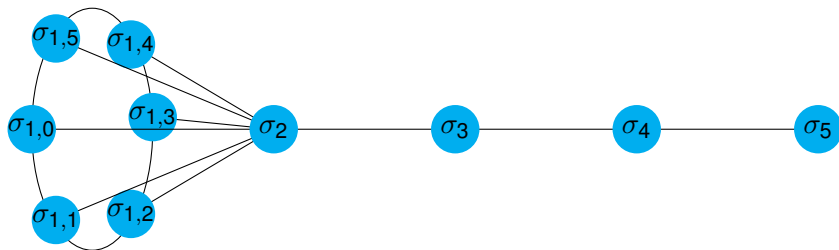
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For this we need to introduce the following remarkable elements, for $1 \leq i \leq j$

- $[j \cdots i] = (j, j-1, \dots, i) = s_{j-1} \cdots s_{i+1} s_i$
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and we set

$$\begin{aligned} [\mathbf{j} \cdots \mathbf{i}] &= \sigma_{j-1} \cdots \sigma_{i+1} \sigma_i \in \mathcal{M} \\ [\mathbf{i} \cdots \mathbf{j}] &= \sigma_i \sigma_{i+1} \cdots \sigma_{j-1} \in \mathcal{M} \end{aligned}$$

where \mathcal{M} is the free monoid on the atoms of B^+ .

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$$w = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & \zeta_7^3 & 0 & 0 \\ 0 & 0 & \zeta_7^4 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix} \in G(7, 7, 4) \rightsquigarrow R_4(w) = [\mathbf{4} \cdots \mathbf{1}] = \sigma_3 \sigma_2 \sigma_1$$

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Then the matrix obtained inside $G(7, 7, 2)$ is equal to s_1 hence

$$R_2(w) = \sigma_1 \text{ and } R(w) = R_2(w) R_3(w) R_4(w) = \sigma_1 \sigma_2 \sigma_{1,4} \sigma_1 \sigma_3 \sigma_2 \sigma_1.$$



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We have $\text{Div}(J) = \mathfrak{S}_n < G(e, e, n)$.

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Lemma ($r = 1$)

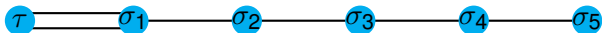
- If \hat{w} is diagonal, then $\ell(s_{1,k} w) = \ell(w) - 1$ if and only if the bottom content of \hat{w} is not 1 ;
- If \hat{w} is antidiagonal, then $\ell(s_{1,k} w) = \ell(w) - 1$ if and only if the top content of \hat{w} is ζ_e^{-k} .

- 1 Definitions
- 2 Standard monoid for $G(e, e, n)$
- 3 Standard monoid for $G(d, 1, n)$
- 4 A word on exceptional groups

Standard monoid for \mathcal{B}_n^*

We consider the monoid B^+ with presentation

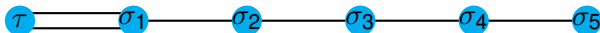
$$\left\langle \tau, \sigma_1, \dots, \sigma_{n-1} \mid \begin{array}{l} \sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1} \\ \sigma_i \sigma_j = \sigma_j \sigma_i, |i - j| \geq 2 \\ \sigma_i \tau = \tau \sigma_i, i > 1 \\ \sigma_1 \tau \sigma_1 \tau = \tau \sigma_1 \tau \sigma_1 \end{array} \right\rangle \quad (1)$$



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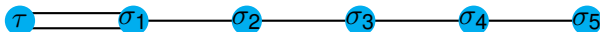
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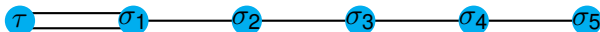
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We denote

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the images of τ and σ_i in W .

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Example

For

$$w = \begin{pmatrix} 0 & 0 & \zeta_7^{-2} \\ \zeta_7^{-1} & 0 & 0 \\ 0 & \zeta_7^2 & 0 \end{pmatrix} \in G(7, 1, 3)$$

one gets

$$\begin{aligned} R_3(w) &= [\mathbf{3} \cdots \mathbf{1}] \tau^5 [\mathbf{1} \cdots \mathbf{2}] = \sigma_2 \sigma_1 \tau^5 \sigma_1 \\ R_2(w) &= \tau \sigma_1 \\ R_1(w) &= \tau^2 \end{aligned}$$

hence

$$R(w) = R_1(w) R_2(w) R_3(w) = \tau^2 \cdot \tau \sigma_1 \cdot \sigma_1 \tau^5 \sigma_1 \sigma_2$$

Neaime's algorithm for \mathcal{B}_n^*

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The set $\text{Div}(c)$ is made of the monomial matrices whose nonzero entries are either 1 or ζ_d^{-1} .

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CRG	Coxeter type	standard monoid
$G(2, 1, n)$	B_n/C_n	classic
$G(2, 2, n)$	D_n	classic
$G(1, 1, n)$	A_{n-1}	classic
$G(e, e, n)$	$I_2(e)$	dual

- 1 Definitions
- 2 Standard monoid for $G(e, e, n)$
- 3 Standard monoid for $G(d, 1, n)$
- 4 A word on exceptional groups**

Definition

The monoid $M(r, s)$ is presented by generators u_1, u_2, \dots, u_r and relations

$$\underbrace{u_1 u_2 u_3 \dots}_s = \underbrace{u_2 u_3 u_4 \dots}_s = \dots = \underbrace{u_r u_1 u_2 \dots}_s$$

where $\underbrace{u_1 u_2 u_3 \dots}_s$ represents the unique subword of length s starting with u_1 of the infinite word $(u_1 u_2 \dots u_r)(u_1 u_2 \dots u_r) \dots$

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$M(r, s)$ is always a Garside monoid, with preferred

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In the cases we are interested in, we shall check that these are actually Garside interval monoids.

- $W = G(e, e, 2)$ with $S = \{s, t\}$,

$$s = (1, 2) \quad t = \begin{pmatrix} 0 & \zeta_e \\ \zeta_e^{-1} & 0 \end{pmatrix}$$

and $c = sts \cdots = tst \cdots$

Then $M(2, e)$ is an interval monoid for B .

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- $W = G_7, G_{11}, G_{19}$ and $G(4, 2, 2) : M(3, 3)$.
- $W = G_{12} : M(3, 4)$
- $W = G_{22} : M(3, 5)$

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- Some of them are *not*, but share the same space X/W as a 'real' group W_0

Let then B^+ the corresponding Artin monoid with set of atoms S . It is an interval monoid w.r.t. W_0 , with preferred Garside element Δ_0 .

Theorem

- *We have $B = \text{Frac}(B^+)$, the atoms of B^+ being mapped to braided reflections.*
- *Setting $|Z(W)| = m$, we have $B^+ = M_S(c)$ where $c = \zeta_m^{-1} \text{Id}$.*

Use of the 'real theory'

W	$ ZW $	$ \text{Div}(c) $	c	W_0	$ ZW_0 $
G_4	2	19	Δ_0^2	$I_2(3)$	1
G_8	4	19	Δ_0^2	$I_2(3)$	1
G_{16}	10	19	Δ_0^2	$I_2(3)$	1
G_5	6	8	Δ_0	$I_2(4)$	2
G_{10}	12	8	Δ_0	$I_2(4)$	2
G_{18}	30	8	Δ_0	$I_2(4)$	2
G_{20}	6	51	Δ_0^2	$I_2(5)$	1
G_6	4	12	Δ_0	$I_2(6)$	2
G_9	8	12	Δ_0	$I_2(6)$	2
G_{17}	20	12	Δ_0	$I_2(6)$	2
G_{21}	12	20	Δ_0	$I_2(10)$	2
G_{25}	3	211	Δ_0^2	A_3	1
G_{26}	6	48	Δ_0	B_3	2
G_{32}	6	3651	Δ_0^2	A_4	1

For the groups G_{24} , G_{27} , G_{29} , G_{33} and G_{34} , none of this works, but Bessis constructed suitable monoids using the fact that they are 'well-generated'¹.

The element c is a Springer regular element, the length is computed from all the reflections, but not all of them are atoms.

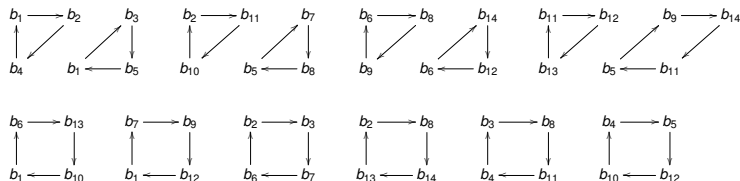
It is then a nontrivial task to get from this a 'short' presentation with geometric meaning.

For the last group G_{31} , one needs a Garside *category* to deal with a corresponding braid *groupoid*, and no Garside monoid is known for this case.

1. applied to the 'real' case, this yields the 'dual braid monoid'

Bessis monoids : example of G_{24}

For $W = G_{24}$, one gets that $M(c)$ is presented by generators b_i , $i = 1, \dots, 14$ such that $b_i = s_i$ for $i \leq 3$, and circular relations depicted as follows



representing the relations

$$\begin{aligned} b_1 b_2 &= b_2 b_4 = b_4 b_1 \\ b_6 b_{13} &= b_{13} b_{10} = b_{10} b_1 = b_1 b_6 \\ \dots \end{aligned}$$