

FUNCTORS AND MORPHISMS DETERMINED BY OBJECTS REVISITED

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- Functors and morphisms determined by objects were introduced in 1978 by Maurice Auslander in his celebrated Philadelphia notes.
- The review begins: *This extremely long paper (244 pages) is devoted to the investigation of functors and morphisms determined by objects.*
- The review ends: *The paper is clearly and concisely written. However, in view of the length of the paper, a table of contents would have been very useful.*

MOTIVATING PROBLEM

- Fix a category, for example a module category.
- Can we classify all **morphisms ending in a fixed object**?
- Two morphisms $\alpha_i: X_i \rightarrow Y$ ($i = 1, 2$) are **isomorphic** if there exists an isomorphism $\phi: X_1 \rightarrow X_2$ such that $\alpha_1 = \alpha_2\phi$.

$$\begin{array}{ccc} X_1 & \xrightarrow{\alpha_1} & Y \\ \downarrow \phi & & \parallel \\ X & \xrightarrow{\alpha_2} & Y \end{array}$$

MORPHISMS DETERMINED BY OBJECTS

DEFINITION (AUSLANDER)

Fix a category \mathcal{C} . A morphism $\alpha: X \rightarrow Y$ in \mathcal{C} is **right determined** by an object C if for every morphism $\alpha': X' \rightarrow Y$ the following are equivalent:

- α' factors through α ;
- for every $\phi: C \rightarrow X'$ the composite $\alpha'\phi$ factors through α .

$$\begin{array}{ccccc} C & \xrightarrow{\phi} & X' & \xrightarrow{\alpha'} & Y \\ & \searrow & \vdots & & \parallel \\ & & X & \xrightarrow{\alpha} & Y \end{array}$$

$\text{Im Hom}(C, \alpha) :=$ image of $\text{Hom}(C, \alpha): \text{Hom}(C, X) \rightarrow \text{Hom}(C, Y)$

The second condition means

$$\text{Im Hom}(C, \alpha') \subseteq \text{Im Hom}(C, \alpha).$$

EXAMPLE: ALMOST SPLIT MORPHISMS

DEFINITION (AUSLANDER–REITEN)

A morphism $\alpha: X \rightarrow Y$ is **right almost split** if α is not a retraction and every morphism $X' \rightarrow Y$ that is not a retraction factors through α .

PROPOSITION (AUSLANDER)

A morphism $\alpha: X \rightarrow Y$ in an additive category is right almost split if and only if

- $\text{End}(Y)$ is a local ring,
- α is right determined by Y ,
- $\text{Im Hom}(Y, \alpha) = \text{rad End}(Y)$.

THE CLASSIFICATION PROBLEM: INVARIANTS

A morphism $\alpha: X \rightarrow Y$ often decomposes as follows: There is a decomposition

$$X = X' \oplus X''$$

such that $\alpha|_{X'}$ is **right minimal** and $\alpha|_{X''} = 0$.

OBSERVATION

Let $\alpha_i: X_i \rightarrow Y$ ($i = 1, 2$) be morphisms that are right minimal and right C -determined. Then

$$\alpha_1 \cong \alpha_2 \quad \iff \quad \text{Im Hom}(C, \alpha_1) = \text{Im Hom}(C, \alpha_2).$$

AUSLANDER'S WORK

Fix a ring Λ and consider the category $\text{Mod } \Lambda$ of Λ -modules.

THEOREM (AUSLANDER)

Let C and Y be Λ -modules with C finitely presented. Given an $\text{End}_\Lambda(C)$ -submodule $H \subseteq \text{Hom}_\Lambda(C, Y)$, there exists a right C -determined morphism $\alpha: X \rightarrow Y$ in $\text{Mod } \Lambda$ satisfying

$$\text{Im Hom}_\Lambda(C, \alpha) = H.$$

THEOREM (AUSLANDER)

Let Λ be an Artin algebra. A morphism $\alpha: X \rightarrow Y$ in $\text{mod } \Lambda$ is right determined as a morphism in $\text{Mod } \Lambda$ by

$$\text{Tr } D(\text{Ker } \alpha) \oplus P(\text{Coker } \alpha).$$

AUSLANDER'S WORK: SOME COMMENTS

- Auslander's formulation of the first result includes a uniqueness statement: α can be chosen to be right minimal.
- The proofs are based on the Auslander–Reiten formula

$$D\underline{\text{Hom}}(X, Y) \cong \text{Ext}^1(Y, D \text{Tr } X).$$

- The first result is a vast generalisation of the existence result for almost split sequences:
- Take for Y a finitely presented module with local endomorphism ring, $C = Y$, and $H = \text{rad } \text{End}(Y)$. Then $\alpha: X \rightarrow Y$ is right almost split.
- The results are not restricted to the category of finitely presented modules; their proofs involve pure-injective modules.

THE CLASSIFICATION PROBLEM: REFORMULATION

Motivated by Auslander's results, we divide the classification problem into two parts:

QUESTION

Fix a category \mathcal{C} and an object $Y \in \mathcal{C}$.

- **Morphisms:** Given an object $C \in \mathcal{C}$ and a subset $H \subseteq \text{Hom}(C, Y)$ which is $\text{End}(C)$ -invariant. Is there a right C -determined morphism $\alpha: X \rightarrow Y$ with $\text{Im Hom}(C, \alpha) = H$?
- **Objects:** Is every morphism ending in Y right determined by an object?

EXAMPLE: PARTIALLY ORDERED SETS

EXAMPLE

Let C be a partially ordered set, viewed as a category.

- Fix a morphism $\alpha: x \rightarrow y$, which means that $x \leq y$.
- If $x = y$, then α is right determined by every object of C .
- If $x \neq y$, then α is right determined by an object $c \in C$ iff

$$C_\alpha = \{c \in C \mid c \not\leq x, c \leq y\}$$

has a unique minimal element. In that case $c = \inf C_\alpha$.

- In (\mathbb{Z}, \leq) , all morphisms are determined by objects.
- In (\mathbb{Q}, \leq) , only identity morphisms are determined by objects.

SERRE FUNCTORS: DEFINITION

DEFINITION (REITEN–VAN DEN BERGH)

Let k be a field and \mathcal{C} a k -linear additive category with finite dimensional Hom-spaces. A **right Serre functor** is an additive functor $S: \mathcal{C} \rightarrow \mathcal{C}$ together with a natural isomorphism

$$D \operatorname{Hom}(X, -) \xrightarrow{\sim} \operatorname{Hom}(-, SX)$$

for all $X \in \mathcal{C}$, where $D = \operatorname{Hom}_k(-, k)$. A right Serre functor is called **Serre functor** if it is an equivalence.

THEOREM

For a Hom-finite triangulated category \mathcal{C} are equivalent:

- *There is a right Serre functor $\mathcal{C} \rightarrow \mathcal{C}$.*
- *Given objects C, Y in \mathcal{C} and an $\text{End}(\mathcal{C})$ -submodule $H \subseteq \text{Hom}(C, Y)$, there exists a morphism $\alpha: X \rightarrow Y$ which is right \mathcal{C} -determined and satisfies $\text{Im Hom}(C, \alpha) = H$.*

Note: A right Serre functor is unique up to isomorphism.

THEOREM

For a right Serre functor $S: \mathcal{C} \rightarrow \mathcal{C}$ are equivalent:

- *The functor S is a Serre functor.*
- *Every morphism in \mathcal{C} is right determined by an object in \mathcal{C} .*

In this case a morphism with cone C is right determined by $S^{-1}C$.

WE THANK HIM:

